An Iterative Algorithm to Compute the Generalized Inverse $A_{T,S}^{(2)}$ Under the Restricted Inner Product

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Abstract—Let T and S be a subspace of C^n and C^m , respectively. Then for $A \in C^{m \times n}$ satisfied $AT \oplus S = C^m$, the generalized inverse $A_{T,S}^{(2)}$ is given by $A_{T,S}^{(2)} = (P_{S^\perp}AP_T)^\dagger$. In this paper, a finite formulae is presented to compute generalized inverse $A_{T,S}^{(2)}$ under the concept of restricted inner product, which defined as $A \cap A > T_{T,S} = P_{S^\perp}AP_T$, $A \cap A > T_{T,S} = P_{T,S}AP_T$. By the initial matrix $A \cap A > T_{T,S}AP_T$, the iterative method, when taken the initial matrix $X_0 = P_T A^* P_{S^{\perp}}$, the generalized inverse $A_{T,S}^{(2)}$ can be obtained within at most mn iteration extens in absence of $T_{T,S}^{(2)}$. steps in absence of roundoff errors. Finally given numerical example is shown that the iterative formulae is quite efficient.

Keywords—Generalized inverse $A_{T,S}^{(2)}$, Restricted inner product, Iterative method, Orthogonal projection.

I. INTRODUCTION

As for our topic, we first recall the definition of generalized inverse $A_{T,S}^{(2)}$ of a matrix $A \in C^{m \times n}$, which has a range Tand null space S.

Definition 1.1^[1]. Let $A \in C^{m \times n}$ be of rank r, T be a subspace of C^n of dimension $s \leq r$ and S be a subspace of C^m of dimension m-s. If X satisfies

$$XAX = X, R(X) = T, N(X) = S$$

then X is called the generalized inverse $A_{T,S}^{(2)}$ of A. It is well-known that the common important generalized inverses are all the generalized inverse $A_{T,S}^{(2)}$, which have the prescribed range T and the null space S of $\{2\}$ -(or outer) inverse of A. These were introduced in [1,2]. The outer generalized inverse $A_{T,S}^{(2)}$ has been widely used in

various fields, for instance, in statistics, control theory, power systems, nonlinear equations, optimization and numerical analysis, and so on. The applications and computations of the outer generalized inverse can be found in [1-12,14-16].

First using finite iterative method to solve the linear system can be seen in [17]. But using this method to compute the generalized inverse only can viewed in author's paper [13]. We could not directly use finite iterative formulae to compute the generalized inverse $A_{T,S}^{(2)}$, because its range and null space are not orthogonal. In this paper, we first define a restricted inner product. Then a restricted norm of a matrix A is generated by this inner product. In the end, a finite iterative formulae to compute the generalized inverse $A_{T,S}^{(2)}$ is devised. Throughout the paper, let $C_r^{m\times n}$ denote the set of all $m\times n$

matrices with rank r over C. T and S is a subspace of C^n

and C^m , respectively, with dimension $s(s \le r)$. For any $A \in$ $C^{m \times n}$, we write R(A) for its range, N(A) for its null range. A^* and r(A) stands for the conjugate transpose and the rank of A, respectively.

Let T and S be the subspace of C^n and C^m , respectively. The restricted conjugate transpose on T and S of $A_{T,S}^*$ for a complex matrix $A \in \hat{C}^{m \times n}$ is defined as $A_{T,S}^* =$ $(P_{S^{\perp}}AP_T)^* = P_TA^*P_{S^{\perp}}$. In the same way, in the space $C^{m \times n}$, an restricted inner product on subspace T and S is defined as $< A, B>_{T,S} = < P_{S^{\perp}}A, BP_T > = tr(A^*P_{S^{\perp}}BP_T)$ for all $A, B \in C^{n \times n}$. Then the restricted norm on the subspace of T and S of a matrix A generated by this inner product is the Frobenius norm of matrix $P_{S^{\perp}}AP_T$ denoted by $||A||_{T,S}$.

In this paper the following Lemmas are needed in what

Lemma 1.1^[1] Let $A \in C^{m \times n}$ be of rank r, any two of the following three statements imply the third:

$$X \in A\{1\}$$

$$X \in A\{2\}$$

$$rankA = rankX$$

Lemma $1.2^{[1]}$ Let $A \in C^{m \times n}$ be of rank r, T be a subspace of C^n of dimension $s \leq r$ and S be a subspace of C^m of dimension m-s. Then A has a $\{2\}$ inverse X such that R(X) = T and N(X) = S, if and only if

$$AT \oplus S = C^m$$

In which case X is unique and it is denoted by $A_{T,S}^{(2)}$.

Lemma $1.3^{[1,2]}$ Let A, T, S and G be the same as in Lemma 1.2. Then A has a $\{2\}$ inverse $A_{T,S}^{(2)}$, and

$$A_{T,S}^{(2)} = (P_{S^{\perp}} A P_T)^{\dagger}$$

Lemma $1.4^{[1]}$ (1) $P_{L,M}A = A$ if and only if $R(A) \subset L$, $(2)AP_{L,M} = A$ if and only if $N(A) \supset M$

Throughout the paper, we assume that $AT \oplus S = C^m(A^*S^\perp \oplus T^\perp = C^n)$, in other words the generalized inverse $A_{T,S}^{(2)}$ of a matrix A is existed.

About the restricted inner product on subspace T and S, we have the following property.

Lemma 1.5 Let T and S be the subspace of C^n and C^m , $A, B \in C^{m \times n}$, then we have:

$$< A, B>_{T,S} = < A, P_{S^{\perp}}BP_{T}> = < P_{S^{\perp}}AP_{T}, B>$$

= $\overline{< B, A>_{T,S}} = < B^{*}, P_{T}A^{*}P_{S^{\perp}}>$

According to the definition and the properties of inner product, the above equalities are right.

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II. Iterative method for computing $A_{T.S}^{(2)}$

In this section we first introduce an iterative method to obtain a solution of the matrix equation $P_{S^\perp}AXAP_T=P_{S^\perp}AP_T$, where $A\in C^{m\times n}$, $T\subset C^n$ and $S\subset C^m$ satisfied $AT\oplus S=C^m$. Then for any initial matrix X_0 with $R(X_0)\subset R(P_TA^*)$ and $N(X_0)\supset N(A^*P_{S^\perp})$, the matrix sequence $\{X_k\}$ generated by the iterative method converges to its a solution within at most mn iteration steps in absence of the roundoff errors. We also show that if let the initial matrix $X_0=P_TA^*P_{S^\perp}$, then the solution X^* obtained by the iterative method is the generalized inverse $A_{T.S}^{(2)}$.

First we present the iteration method for solving the matrix equation $P_{S^{\perp}}AXAP_T = P_{S^{\perp}}AP_T$, the iteration method as follow:

Algorithm 2.1

1. In put matrices $A \in C^{m \times n}$, $P_{S^{\perp}} \in C^{m \times m}$, $P_T \in C^{n \times n}$ and $X_0 \in C^{n \times m}$ with $R(X_0) \subset R(P_T A^*)$ and $N(X_0) \supset N(A^* P_{S^{\perp}})$;

2. Calculate

$$R_0 = A - AX_0A;$$

 $P_0 = A(R_0)^*_{T,S}A;$
 $k := 0:$

3. If $P_{S^{\perp}}R_kP_T=0$, then stop; else, k:=k+1;

4. Calculate

$$\begin{split} X_k &= X_{k-1} + \frac{\parallel R_{k-1} \parallel_{T,S}^2}{\parallel P_{k-1} \parallel_{T,S}^2} (P_{k-1})_{T,S}^*; \\ R_k &= A - AX_k A \\ &= R_{k-1} - \frac{\parallel R_{k-1} \parallel_{T,S}^2}{\parallel P_{k-1} \parallel_{T,S}^2} A(P_{k-1})_{T,S}^* A; \\ P_k &= A(R_k)_{T,S}^* A + \frac{\parallel R_k \parallel_{T,S}^2}{\parallel R_{k-1} \parallel_{T,S}^2} P_{k-1}; \end{split}$$

5. Goto step3.

About Algorithm 2.1, we have the following basic properties.

Theorem 2.2 In Algorithm 2.1, if we take the initial matrix $X_0 = A_{T,S}^*$, then the sequences $\{X_i\}$ and $\{P_i\}$ generalized by it such that

 $\begin{array}{cccc} (1)R(X_k) &\subset R(P_TA^*), & N(X_k) \supset N(A^*P_{S^\perp}) & \text{and} \\ R(P_k) \subset R(AP_T), & N(P_k) \supset N(P_{S^\perp}A); & \end{array}$

(2)if $AT \oplus S = C^m$ and $P_{S^{\perp}}R_kP_T = 0$, then $X_k = A_{T,S}^{(2)}$; **Proof** (1)To prove the conclusion, we use the induction When i = 0, we have

$$\begin{array}{rcl} X_0 & = & A_{T,S}^* = P_T A^* P_{S^{\perp}} \\ P_0 & = & A(R_0)_{T,S}^* A = A P_T R_0^* P_{S^{\perp}} A \end{array}$$

this implies the conclusion is right.

When i = 1, we have

$$X_{1} = X_{0} + \frac{\parallel R_{0} \parallel_{T,S}^{2}}{\parallel P_{0} \parallel_{T,S}^{2}} (P_{0})_{T,S}^{*}$$

$$= P_{T}A^{*}(P_{S^{\perp}} + \frac{\parallel R_{0} \parallel_{T,S}^{2}}{\parallel P_{0} \parallel_{T,S}^{2}} P_{S^{\perp}}R_{0}P_{T}A^{*}P_{S^{\perp}})$$

$$= (P_{T} + \frac{\parallel R_{0} \parallel_{T,S}^{2}}{\parallel P_{0} \parallel_{T,S}^{2}}) P_{T}A^{*}P_{S^{\perp}}R_{0}P_{T})A^{*}P_{S^{\perp}}$$

and

$$P_{1} = AP_{T}R_{1}^{*}P_{S^{\perp}}A + \frac{\parallel R_{1} \parallel_{T,S}^{2}}{\parallel R_{0} \parallel_{T,S}^{2}}P_{0}$$

$$= AP_{T}(P_{T}R_{1}^{*}P_{S^{\perp}}A + \frac{\parallel R_{1} \parallel_{T,S}^{2}}{\parallel R_{0} \parallel_{T,S}^{2}}P_{T}R_{0}^{*}P_{S^{\perp}}A)$$

$$= (AP_{T}R_{1}^{*}P_{S^{\perp}} + \frac{\parallel R_{1} \parallel_{T,S}^{2}}{\parallel R_{0} \parallel_{T,S}^{2}}AP_{T}R_{0}^{*}P_{S^{\perp}})P_{S^{\perp}}A$$

This shows when i = 1, the conclusion is also right.

Assume that conclusion holds for all $0 \le i \le s(0 < s < k)$. Then there exist matrices U, V, W, and Y such that

$$X_s = P_T A^* U = V A^* P_{S^{\perp}}$$

 $P_s = A P_T W = Y P_{S^{\perp}} A$

Further, we have that

$$X_{s+1} = X_s + \frac{\parallel R_s \parallel_{T,S}^2}{\parallel P_s \parallel_{T,S}^2} (P_s)_{T,S}^*$$

$$= P_T A^* U + \frac{\parallel R_s \parallel_{T,S}^2}{\parallel P_s \parallel_{T,S}^2} P_T A^* P_{S^{\perp}} Y^* P_{S^{\perp}}$$

$$= P_T A^* (U + \frac{\parallel R_s \parallel_{T,S}^2}{\parallel P_s \parallel_{T,S}^2} P_{S^{\perp}} Y^* P_{S^{\perp}})$$

$$= V A^* P_{S^{\perp}} + \frac{\parallel R_s \parallel_{T,S}^2}{\parallel P_s \parallel_{T,S}^2} P_T W^* P_T A^* P_{S^{\perp}}$$

$$= (V + \frac{\parallel R_s \parallel_{T,S}^2}{\parallel P_s \parallel_{T,S}^2} P_T W^* P_T) A^* P_{S^{\perp}}$$

and

$$P_{s+1} = AP_{T}R_{s+1}^{*}P_{S^{\perp}}A + \frac{\parallel R_{s+1} \parallel_{L}^{2}T, S}{\parallel R_{s} \parallel_{T,S}^{2}}P_{s}$$

$$= AP_{T}(R_{s+1}^{*}P_{S^{\perp}}A + \frac{\parallel R_{s+1} \parallel_{T,S}^{2}}{\parallel R_{s} \parallel_{T,S}^{2}}W)$$

$$= (AP_{T}R_{s+1}^{*} + \frac{\parallel R_{s+1} \parallel_{T,S}^{2}}{\parallel R_{s} \parallel_{T,S}^{2}}Y)P_{S^{\perp}}A$$

This implies that

$$R(X_{s+1}) \subset R(P_LA^*P_L) \quad \text{and} \quad N(X_{s+1}) \supset N(P_LA^*P_L)$$
 and

$$R(P_{s+1}) \subset R(AP_L)$$
 and $N(P_{s+1}) \supset N(P_LA)$

By the principle of induction, the conclusion holds for all $i=0,1,\cdots$

(2)From the lemma 1.4 and the results in (1) of the Theorem 2.2, we can easy to get $P_TX=XP_{S^\perp}=X$. According to Algorithm 2.1, we know that, if $P_{S^\perp}R_kP_T=0$, then we have $X_k\in (P_{S^\perp}AP_T)\{1\}$. This implies $r(X_k)\geq r(P_{S^\perp}AP_T)$, then by the conclusion of (1) and $AT\oplus S=C^m$, we can easy get $r(X_k)=r(P_{S^\perp}AP_T)$. From Lemma 2.1 we know $X_k\in (P_{S^\perp}AP_T)\{1,2\}$ with range $R(P_TA^*P_{S^\perp})$ and null space $N(P_TA^*P_{S^\perp})$. If $AT\oplus S=C^m$, then $R(P_TA^*P_{S^\perp})=R(P_TA^*)=T$ and $N(P_TA^*P_{S^\perp})=N(A^*P_{S^\perp})=S$. By Lemma 1.3, it shows that $X_k=(P_{S^\perp}AP_T)^\dagger=A_{T,S}^{(2)}$.

Theorem 2.3 Let \widetilde{X} be an solution of matrix equation $P_{S^{\perp}}AXAP_T = P_{S^{\perp}}AP_T$ with $R(X) \subset T$ and $N(X) \subset T$ S, then for any initial matrix X_0 with $R(X_0) \subset T$ and $N(X_0) \subset S$, the sequences $\{X_i\}$, $\{R_i\}$ and $\{P_i\}$ generalized by Algorithm 3.1 satisfy

$$< P_i, P_{S^{\perp}}(\widetilde{X} - X_i)^* P_T >_{T,S} = ||R_i||_{T,S}^2, (i = 0, 1, 2, \cdots).$$

Proof We prove the conclusion by induction. By Algorithm 2.1 and lemma 1.4, when i = 0, we have

$$< P_{0}, P_{S^{\perp}}(\widetilde{X} - X_{0})^{*}P_{T} >_{T,S}$$

$$= < P_{S^{\perp}}P_{0}P_{T}, P_{S^{\perp}}(\widetilde{X} - X_{0})^{*}P_{T} >$$

$$= < P_{0}, P_{S^{\perp}}(\widetilde{X} - X_{0})^{*}P_{T} >$$

$$= < AP_{T}R_{0}^{*}P_{S^{\perp}}A, (\widetilde{X} - X_{0})^{*} >$$

$$= < P_{T}R_{0}^{*}P_{S^{\perp}}, A^{*}(\widetilde{X} - X_{0})^{*}A^{*} >$$

$$= < P_{T}R_{0}^{*}P_{S^{\perp}}, R_{0}^{*} > = ||R_{0}||_{T,S}^{2}$$

And when i = 1, we have

$$< P_{1}, P_{S^{\perp}}(\widetilde{X} - X_{1})^{*}P_{T} >_{T,S}$$

$$= < P_{S^{\perp}}P_{1}P_{T}, P_{S^{\perp}}(\widetilde{X} - X_{1})^{*}P_{T} >$$

$$= < P_{1}, P_{S^{\perp}}(\widetilde{X} - X_{1})^{*}P_{T} >$$

$$= < P_{1}, (\widetilde{X} - X_{1})^{*} > \text{ by Theorem 2.2}$$

$$= \left< AP_{T}R_{1}^{*}P_{S^{\perp}}A + \frac{\parallel R_{1} \parallel_{T,S}^{2}}{\parallel R_{0} \parallel_{T,S}^{2}}P_{0}, (\widetilde{X} - X_{1})^{*} \right>$$

$$= < AP_{T}R_{1}^{*}P_{S^{\perp}}A, (\widetilde{X} - X_{1})^{*} >$$

$$+ \frac{\parallel R_{1} \parallel_{T,S}^{2}}{\parallel R_{0} \parallel_{T,S}^{2}} < P_{0}, (\widetilde{X} - X_{1})^{*} >$$

$$= < P_{T}R_{1}^{*}P_{S^{\perp}}, R_{1}^{*} > + \frac{\parallel R_{1} \parallel_{T,S}^{2}}{\parallel R_{0} \parallel_{T,S}^{2}} < P_{0}, (\widetilde{X} - X_{0})^{*} >$$

$$- \frac{\parallel R_{1} \parallel_{T,S}^{2}}{\parallel P_{0} \parallel_{T,S}^{2}} < P_{0}, (P_{T}P_{0}^{*}P_{S^{\perp}})^{*} >$$

$$= \parallel R_{1} \parallel_{T,S}^{2}$$

$$= \parallel R_{1} \parallel_{T,S}^{2}$$

Assume that the conclusion holds for i = s(s > 0), that $< P_s, P_{S^{\perp}}(X - X_s)^* P_T >_{T,S} = ||R_s||_{T,S}^2$, then i = s + 1,

$$\langle P_{s+1}, P_{S^{\perp}}(\widetilde{X} - X_{s+1})^* P_T \rangle_{T,S}$$
 and
$$= \langle P_{S^{\perp}} P_{s+1} P_T, P_{S^{\perp}}(\widetilde{X} - X_{s+1})^* P_T \rangle$$

$$= \langle P_{s+1}, P_{S^{\perp}}(\widetilde{X} - X_{s+1})^* P_T \rangle$$

$$= \langle P_{s+1}, (\widetilde{X} - X_{s+1})^* \rangle$$
 by Theorem 2.2
$$= \left\langle A P_T R_{s+1}^* P_{S^{\perp}} A + \frac{\parallel R_{s+1} \parallel_{T,S}^2}{\parallel R_s \parallel_{T,S}^2} P_s, (\widetilde{X} - X_{s+1})^* \right\rangle$$

$$= \langle A P_T R_{s+1}^* P_{S^{\perp}} A, (\widetilde{X} - X_{s+1})^* \rangle$$

$$+ \frac{\parallel R_{s+1} \parallel_{T,S}^2}{\parallel R_s \parallel_{T,S}^2} \langle P_s, (\widetilde{X} - X_{s+1})^* \rangle$$

$$= \langle P_T R_{s+1}^* P_{S^{\perp}}, R_{s+1}^* \rangle + \frac{\parallel R_{s+1} \parallel_{T,S}^2}{\parallel R_s \parallel_{T,S}^2} \langle P_s, (\widetilde{X} - X_s)^* \rangle$$

$$- \frac{\parallel R_{s+1} \parallel_{T,S}^2}{\parallel P_s \parallel_{T,S}^2} \langle P_s, P_{S^{\perp}} P_s P_T \rangle$$

$$= \| R_{s+1} \|_{T,S}^2$$

$$= \| R_{s+1} \|_{T,S}^2$$

By the principle of induction, the conclusion $< P_i, P_{S^\perp}(\widetilde{X} - X_i)^*P_T>_{T,S} = \parallel R_i \parallel_{T,S}^2$ holds for all

Remark 1 From Theorem 2.3 we know that if $P_{S^{\perp}}R_iP_T \neq$ 0, then $P_{S^{\perp}}P_iP_T \neq 0$. This result shows that if $P_{S^{\perp}}R_iP_T \neq 0$ 0, then Algorithm 3.1 can not be terminated.

Theorem 2.4 For the sequences $\{R_i\}$ and $\{P_i\}$ generated by Algorithm 2.1 with the $X_0 = P_T A^* P_{S^{\perp}}$, if there exists a positive number k such that $R_i \neq 0$ for all $i = 0, 1, 2, \dots k$, then we have

$$\langle R_i, R_j \rangle_{T,S} = 0, \langle P_i, P_j \rangle_{T,S} = 0, (i \neq j, i, j = 0, 1, \dots, k)$$

Proof According to Lemma 1.5, we know that

$$\langle A, B \rangle_{TS} = \overline{\langle B, A \rangle_{TS}}$$

holds for all matrices A and B in $C^{m \times n}$, so we only need prove the conclusion hold for all $0 \le i < j \le k$. Using induction and two steps are required.

Step1. Show that $\langle R_i, R_{i+1} \rangle_{T,S} = 0$ and \langle $P_i, P_{i+1} >_{T,S} = 0$ for all $i = 0, 1, 2, \dots, k$. To prove this conclusion, we also use induction. According to Lemma 1.5 and Algorithm 2.1, when i = 0, we have

$$< R_{0}, R_{1} >_{T,S} = < P_{S^{\perp}} R_{0} P_{T}, R_{1} >$$

$$= \left\langle P_{S^{\perp}} R_{0} P_{T}, R_{0} - \frac{\parallel R_{0} \parallel_{T,S}^{2}}{\parallel P_{0} \parallel_{T,S}^{2}} A P_{T} P_{0}^{*} P_{S^{\perp}} A \right\rangle$$

$$= < P_{S^{\perp}} R_{0} P_{T}, R_{0} > -\frac{\parallel R_{0} \parallel_{T,S}^{2}}{\parallel P_{0} \parallel_{T,S}^{2}} < P_{S^{\perp}} R_{0} P_{T}, A P_{T} P_{0}^{*} P_{S^{\perp}} A >$$

$$= \parallel R_{0} \parallel_{T,S}^{2} - \frac{\parallel R_{0} \parallel_{T,S}^{2}}{\parallel P_{0} \parallel_{T,S}^{2}} < A^{*} P_{S^{\perp}} R_{0} P_{T} A^{*}, P_{T} P_{0}^{*} P_{S^{\perp}} >$$

$$= \parallel R_{0} \parallel_{T,S}^{2} - \frac{\parallel R_{0} \parallel_{T,S}^{2}}{\parallel P_{0} \parallel_{T,S}^{2}} < P_{0}^{*}, P_{T} P_{0}^{*} P_{S^{\perp}} >$$

$$= \parallel R_{0} \parallel_{T,S}^{2} - \frac{\parallel R_{0} \parallel_{T,S}^{2}}{\parallel P_{0} \parallel_{T,S}^{2}} \parallel P_{0} \parallel_{T,S}^{2} = 0$$

and

$$= \langle P_{S^{\perp}}P_{s+1}P_{T}, P_{S^{\perp}}(X - X_{s+1})^{*}P_{T} \rangle$$

$$= \langle P_{s+1}, P_{S^{\perp}}(\widetilde{X} - X_{s+1})^{*}P_{T} \rangle$$

$$= \langle P_{s+1}, (\widetilde{X} - X_{s+1})^{*} \rangle$$

$$= \langle P_{s+1}P_{s+1}P_{s+1}A + \frac{\|R_{s+1}\|_{T,S}^{2}}{\|R_{s}\|_{T,S}^{2}} P_{s}, (\widetilde{X} - X_{s+1})^{*} \rangle$$

$$= \langle AP_{T}R_{s+1}^{*}P_{S^{\perp}}A + \frac{\|R_{s+1}\|_{T,S}^{2}}{\|R_{s}\|_{T,S}^{2}} \langle P_{s}, (\widetilde{X} - X_{s+1})^{*} \rangle$$

$$= \langle AP_{T}R_{s+1}^{*}P_{S^{\perp}}A, (\widetilde{X} - X_{s+1})^{*} \rangle$$

$$= \langle AP_$$

Assume that conclusion holds for all $i \le s(0 < s < k)$,

then

and

$$< P_{s}, P_{s+1} >_{T,S} = < P_{S^{\perp}} P_{s} P_{T}, P_{s+1} >$$

$$= \left\langle P_{S^{\perp}} P_{s} P_{T}, A P_{T} R_{s+1}^{*} P_{S^{\perp}} A + \frac{\parallel R_{s+1} \parallel_{T,S}^{2}}{\parallel R_{s} \parallel_{T,S}^{2}} P_{s} \right\rangle$$

$$= < A^{*} P_{S^{\perp}} P_{s} P_{T} A^{*}, P_{T} R_{s+1}^{*} P_{S^{\perp}} > +$$

$$\frac{\parallel R_{s+1} \parallel_{T,S}^{2}}{\parallel R_{s} \parallel_{T,S}^{2}} < P_{S^{\perp}} P_{s} P_{T}, P_{s} >$$

$$= \frac{\parallel P_{s} \parallel_{T,S}^{2}}{\parallel R_{s} \parallel_{T,S}^{2}} < (R_{s} - R_{s+1})^{*}, P_{T} R_{s+1}^{*} P_{S^{\perp}} > +$$

$$\frac{\parallel R_{s+1} \parallel_{T,S}^{2}}{\parallel R_{s} \parallel_{T,S}^{2}} \parallel P_{s} \parallel_{T,S}^{2}$$

$$= -\frac{\parallel P_{s} \parallel_{T,S}^{2}}{\parallel R_{s} \parallel_{T,S}^{2}} \parallel R_{s+1} \parallel_{T,S}^{2} + \frac{\parallel R_{s+1} \parallel_{T,S}^{2}}{\parallel R_{s} \parallel_{T,S}^{2}} \parallel P_{s} \parallel_{T,S}^{2} = 0$$

By the principle of induction, $\langle R_i, R_{i+1} \rangle_{T,S} = 0$, and $\langle P_i, P_{i+1} \rangle_{T,S} = 0$, hold for all $i = 0, 1, \dots, k$.

Step2. Assume that $< R_i, R_{i+l} >_{T,S} = 0$, and $< P_i, P_{i+l} >_{T,S} = 0$, hold for all $0 \le i \le k$ and 1 < l < k, show that $< R_i, R_{i+l+1} >_{T,S} = 0$, and $< P_i, P_{i+l+1} >_{T,S} = 0$.

$$< R_{i}, R_{i+l+1} >_{T,S} = < P_{S^{\perp}} R_{i} P_{T}, R_{i+l+1} >$$

$$= \left\langle P_{S^{\perp}} R_{i} P_{T}, R_{i+l} - \frac{\parallel R_{i+l} \parallel_{T,S}^{2}}{\parallel P_{i+l} \parallel_{T,S}^{2}} A P_{T} P_{i+l}^{*} P_{S^{\perp}} A \right\rangle$$

$$= -\frac{\parallel R_{i+l} \parallel_{T,S}^{2}}{\parallel P_{i+l} \parallel_{T,S}^{2}} < P_{S^{\perp}} R_{i} P_{T}, A P_{T} P_{i+l}^{*} P_{S^{\perp}} A >$$

$$= -\frac{\parallel R_{i+l} \parallel_{T,S}^{2}}{\parallel P_{i+l} \parallel_{T,S}^{2}} < A^{*} P_{S^{\perp}} R_{i} P_{T} A^{*}, P_{T} P_{i+l}^{*} P_{S^{\perp}} >$$

$$= -\frac{\parallel R_{i+l} \parallel_{T,S}^{2}}{\parallel P_{i+l} \parallel_{T,S}^{2}} \left\langle P_{i} - \frac{\parallel R_{i} \parallel_{T,S}^{2}}{\parallel R_{i-1} \parallel_{T,S}^{2}} P_{i-1}, P_{T} P_{i+l}^{*} P_{S^{\perp}} \right\rangle$$

$$= 0$$

and

From step 1 and step 2, we have by principle induction that $\langle R_i, R_j \rangle_{T,S} = 0$, and $\langle P_i, P_j \rangle_{T,S} = 0$, hold for all $i, j = 0, 1, \dots, k, i \neq j$.

Remark 2 Theorem 2.4 implies that, for an initial matrix $X_0 = P_T A^* P_{S^\perp}$, since the R_0, R_1, \cdots are orthogonal each other, based on restricted inner product on subspace T and S, in the finite dimension matrix space $C^{m \times n}$, it is certainly there exists a positive number $k \leq mn$ such that $\|R_k\|_{T,S} = 0$. Then by Theorem 2.2, the generalized inverse $A_{T,S}^{(2)}$ can be obtained within at most mn iteration steps.

Remark 3 When $T=R(A^*)$ and $S=N(A^*)$, the Algorithm 2.1 is exact the finite iterative method to computed M-P inverse A^{\dagger} , which can be reviewed in [13]. Here we recalled it as following:

Algorithm $2.2^{[13]}$

- 1. In put matrices $A \in C^{m \times n}$ and $X_0 \in C^{n \times m} = A^*$;
- 2. Calculate

$$R_0 = A - AX_0A;$$

$$P_0 = AR_0^*A;$$

$$k := 0.$$

- 3. If $R_k = 0$, then stop; else, k := k + 1;
- 4. Calculate

$$\begin{split} X_k &= X_{k-1} + \frac{\parallel R_{k-1} \parallel^2}{\parallel P_{k-1} \parallel^2} P_{k-1}^*; \\ R_k &= A - AX_k A \\ &= R_{k-1} - \frac{\parallel R_{k-1} \parallel^2}{\parallel P_{k-1} \parallel^2} A P_{k-1}^* A; \\ P_k &= AR_k^* A + \frac{\parallel R_k \parallel^2}{\parallel R_{k-1} \parallel^2} P_{k-1}; \end{split}$$

5. Goto step 3.

Remark 4 If $G \in C^{n \times m}$ satisfied R(G) = T and N(G) = S with m,n are large, we can use the Algorithm 2.2 to compute the G^{\dagger} , the $P_T = P_{R(G)} = GG^{\dagger}$ and $P_{S^{\perp}} = P_{N(G)^{\perp}} = G^{\dagger}G$.

III. NUMERICAL EXAMPLES

In this section, we will give a numerical example ,which is taken from [11], to illustrate our results. All the tests are performed by MATLAB 6.1 and the initial iterative matrices are chosen as $X_0 = P_T A^* P_{S^\perp}$. Because of the influence of the error of roundoff, we regard the matrix $P_{S^\perp} A P_T$ as zero matrix if $\parallel A \parallel_{T,S} < 10^{-10}$.

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Example 3.1 Take

$$A = \begin{pmatrix} -1 & 2 & 1 & 0 \\ 1 & 0 & 1 & 1 \\ -1 & -3 & 1 & 2 \end{pmatrix} \in R_3^{3 \times 4},$$

$$G = \begin{pmatrix} 3 & 1 & 0 \\ -2 & 4 & -2 \\ -5 & -4 & 1 \\ 0 & 7 & -3 \end{pmatrix} \in R_2^{4 \times 3}$$

We can easy show that $AR(G)\oplus N(G)=R^3$, then by Lemma 2.1 we know that $A^{(2)}_{R(G),N(G)}$ exists.

By directly computing,

$$P_{R(G)} = \frac{1}{59} \begin{pmatrix} 14 & -10 & -23 & -1 \\ -10 & 24 & 8 & 26 \\ -23 & 8 & 42 & -11 \\ -1 & 26 & -11 & 38 \end{pmatrix}$$

and

$$P_{N(G)^{\perp}} = P_{R(G^T)} = \frac{1}{59} \begin{pmatrix} 102 & 13 & 9\\ 13 & 30 & -11\\ 9 & -11 & 6 \end{pmatrix}$$

Using Algorithm 2.1 and iterate 11 steps, we have X_{11} as follow:

$$X_{11} = \left(\begin{array}{cccc} -0.27419354838934 & 0.32258064516088 & -0.17741935483885 \\ 0.09677419355105 & 0.70967741935533 & -0.29032258064499 \\ 0.50000000000316 & -0.9999999999942 & 0.5000000000020 \\ -0.12903225806276 & 1.38709677419387 & -0.61290322580634 \end{array} \right)$$

with

$$||R_{11}||_{T,S}^2 = ||A - AX_{11}A||_{T,S}^2 = 1.950379437652951 \times 10^{-21}$$

On other hand, by computing, we obtain that

$$A_{T,S}^{(2)} = \frac{1}{62} \begin{pmatrix} -17 & 20 & -11 \\ 6 & 44 & -18 \\ 31 & -62 & 31 \\ -8 & 86 & -38 \end{pmatrix}$$

Then from the above data, we can find that the iterative sequence $\{X_k\}$ converges to $A_{T,S}^{(2)}$.

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