# The More Organized Proof For Acyclic Coloring Of Graphs With $\Delta=5$ with 8 Colors 

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#### Abstract

An acyclic coloring of a graph $G$ is a coloring of its vertices such that:(i) no two neighbors in $G$ are assigned the same color and (ii) no bicolored cycle can exist in $G$. The acyclic chromatic number of $G$ is the least number of colors necessary to acyclically color $G$. Recently it has been proved that any graph of maximum degree 5 has an acyclic chromatic number at most 8 . In this paper we present another proof for this result.


## I. Introduction

AProper coloring of a graph $G$ is a coloring of its vertices such that no two neighbors in $G$ are assigned the same color. An acyclic coloring of a graph $G$ is a proper coloring such that the graph induced by two colors $\alpha$ and $\beta$ is a forest. The minimum number of colors necessary to acyclically color $G$ is called the acyclic chromatic number of $G$ and denoted by $a(G)$. For a family $F$ of graphs, the acyclic chromatic number of $F$, denoted by $a(F)$ is defined as follow : $a(F)=$ $\max \{a(G)$ forall $G \in F\}$.
$a(F)$ has been determined for several families of graphs such as planar graphs [4], 1-planar graphs [2],planar graphs with large girth [3],outer planar graphs [11], product of trees [9], the graphs with maximum degree 3 [8], [10], the graphs with $\Delta=4$ [5], Alon et al [1] showed that
(1):Asymptotically there exist graphs of maximum degree $\Delta$ with acyclic chromatic number in $\Omega\left(\frac{\Delta^{4 / 3}}{\left(\log \Delta^{1 / 3}\right.}\right)$.
(2): Asymptotically it is possible to acyclically color any graph of maximum degree $\Delta$ with $O\left(\Delta^{4 / 3}\right)$ colors.
(3):Trivial greedy polynomial time algorithm exists that acycliccally colors any graphs of maximum degree $\Delta$ with $\Delta^{2}+1$ colors. Fertin and Raspaud [7] proved that nine colors are enough for acyclic coloring a graph with $\Delta=5$. Kishore Yadav etal [12] showed that any graph with $\Delta=5$ can be acyclically colored with 8 colors. In this paper we achieve the above result by another approach which is easier than what has been presented before.

## II. Preliminaries

In the following we only consider graphs of maximum degree $\Delta=5$. Let $N(u)$ be the neighbors set of vertex $u$ and $c(u)$ denoted the color of $u$. The set of colors are assigned to vertices in $N(u)$, denoted by $S C N(u)$. The color $\alpha \in\{1,2,3,4,5,6,7,8\}$ is regarded as a free color for $u$ when:

1) If no color assigned to $u$, then $\alpha \notin S C N(u)$.

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2) If $u$ is colored, then $\alpha \notin S C N(u) \cup\{c(u)\}$.

The set of free colors for $u$ is named $f(u)$. The number of different colors in the neighbors of $u$ is denoted by $d c n(u)$ and we define the color list as follow:
$L_{u}=\left(n_{1}, n_{2}, \ldots, n_{d c n(u)}\right)$ (where $n_{1} \geq n_{2} \geq \ldots \geq n_{d c n(u)}$ ) where an $n_{i}$ represents for a color $\alpha$ in $S C N(u)$ and it is the number of times , $\alpha$ is used among the colored neighbors of $u$.
The color $\alpha$ is free for $u$ if $\alpha \in f(u)$, and $\alpha$ is a valid color for $u$ if $1 \leq \alpha \leq 8$ and assigning it to $u$, still results in an acycliclly coloring. Let $\alpha, \beta$ are two distinct colors. A critical cycle denoted by $C_{u}(\alpha, \beta)$, is a cycle such as $C$, involving $u$ in which all vertices in $C$ are alternatively colored by $\alpha$ and $\beta$ moreover $c(u) \notin\{\alpha, \beta\}$. We can't assign the color $\alpha$ and $\beta$ to $u$ (since they are not valid colors). A vertex $u$ is called single if all its neighbors receive distinct colors.

## III. Acyclic Coloring of Graphs of Maximum Degree 5

In this section we show that 8 color is enough for acyclically coloring of any graph with $\Delta=5$. At first we prove 4 lemmas and in Theorem 1 we will achieve the goal.

Lemma 3.1: If $u$ is an uncolored vertex and $L_{u}=$ $(1,1,1,1,1)$ then we can color $u$ with a valid color

Proof: Since $L_{u}=(1,1,1,1,1)$ so $f(u) \geq 3$ and we can assign one of the color in $f(u)$ such that $1 \leq c(u) \leq 8$.
Lemma 3.2: If $u$ is an uncolored vertex and $L_{u}=$ $(2,1,1,1)$ and $9 \notin S C N(u)$, then we can find a valid color for $u$.

Proof: Let $N(u)=\{v, w, x, y, z\}$ and $c(v)=c(w)=$ $1, c(x)=2, c(y)=3$ and $c(z)=4$,therefore $f(u)=$ $\{5,6,7,8\}$. If we can't choose a valid color from $f(u)$, then we must have $C_{u}(1,5), C_{u}(1,6), C_{u}(1,7)$ and $C_{u}(1,8)$. This means that $v$ and $w$ are single vertices. By assigning color 1 to $u$ and eliminating the colors of $v$ and $w$, we will have $L_{v}=(1,1,1,1,1), L_{w}=(1,1,1,1,1)$ and by Lemma 3.1 we can color $v$ and $w$ with a valid color.
Lemma 3.3: if $u$ is a colored vertex with a valid color and $L_{u}=(2,1,1,1), 9 \notin S C N(u)$. Then we can recolor $u$ with a valid color.

Proof: Let $N(u)=\{v, w, x, y, z\}$ and $c(v)=c(w)=$ $1, c(x)=2, c(y)=3, c(z)=4$ and $c(u)=5$,therefore $f(u)=$ $\{6,7,8\}$. If we can't find a valid color from $f(u)$ to change the color of $u$,then we must have $C_{u}(1,6), C_{u}(1,7)$ and $C_{u}(1,8)$. Now by eliminating the colors of $v$ and $w$ and assigning color 1 to $u$ we have $S C N(v)=\{1,6,7,8, \alpha\}$ and $S C N(w)=$ $\{1,6,7,8, \beta\}$. Let us detail the possible cases for $\alpha$ :
3.3.1 If $\alpha \notin\{6,7,8\}$ then $L_{v}=(1,1,1,1,1)$ and with Lemma 3.1 we can color $v$ with a valid color.
3.3.2 If $\alpha \in\{6,7,8\}$ then $L_{v}=(2,1,1,1)$ and by Lemma 3.2 we can color vertex $v$ with a valid color. We have similar cases for $\beta$ and we can color $w$ with a valid color.
Lemma 3.4: if $u$ is a colored with a valid color with $L_{v}=$ $(2,1,1,1)$ and one of its neighbors is colored with color 9 , then we can recolor vertex $u$ with a valid color.

Proof: Let $N(u)=\{v, w, x, y, z\}$ and $c(v)=c(w)=$ $1, c(x)=2, c(y)=3, c(z)=9$ and $c(u)=4$ therefore $f(u)=$ $\{5,6,7,8\}$. If we can't find a valid color from $f(u)$ to vertex $u$ then we must have $C_{u}(1,5), C_{u}(1,6), C_{u}(1,7)$ and $C_{u}(1,8)$. In this case by eliminating the colors of $v$ and $w$ and assigning color 1 to vertex $u$, we will have $L_{v}=L_{w}=(1,1,1,1,1)$ and by Lemma 3.1 we can find valid colors for $v$ and $w$. To have acyclic coloring of a graph $G$ with 8 colors, first we add $5-d(u)$ new vertices (vertex) for every vertex $u \in V(G)$ and insert edges between $u$ and these new vertices. By the above operation we get a new graph $G^{\prime}$ with the following properties:

$$
d(u)= \begin{cases}5 & \text { if } u \in V(G)  \tag{1}\\ 1 & \text { otherwise }\end{cases}
$$

If we use the algorithm Fertin and Raspaud [7] for graph $G^{\prime}$, then $G^{\prime}$ can be colored acyclically with 9 colors. Then we try to recolor every vertex of $G^{\prime}$ that its color is 9 with a valid color. Finally by removing all vertices of degree 1 we achieve the goal.

Theorem 3.1: Let $G^{\prime}$ is a graph with maximum degree 5 and acyclically colored with 9 colors and let $u$ be a vertex such that $d(u)=5$ and $c(u)=9$ then we can find a valid color to $u$.

Proof: Let us detail possible cases:
Case 3.1.1: $L_{u}=(1,1,1,1,1)$
In this case by eliminating the color of $u$ and using
Lemma 3.1, we can find a valid color for $u$.
Case 3.1.2: $L_{u}=(2,1,1,1)$
In this case by eliminating the color of $u$ and using
Lemma 3.2, we can find a valid color for vertex $u$.
Case 3.1.3: $L_{u}=(3,1,1)$
Let $N(u)=\{v, w, x, y, z\}$ and $c(v)=c(w)=$ $c(x)=1, c(y)=2, c(z)=3$ and $c(u)=9$. These assumptions imply that $f(u)=\{4,5,6,7,8\}$. If we can't choose a color from $f(u)$ to recolor $u$,then we must have $C_{u}(1,4), C_{u}(1,5), C_{u}(1,6), C_{u}(1,7)$ and $C_{u}(1,8)$. The above discussion shows that $N(v) \cup$ $N(w) \cup N(x)-\{u\}$ contains two vertices of color 4 ,two vertices of color 5 ,two vertices of color 6 ,two vertices of color 7 and two vertices of color 8 . So we have at least 10 vertices in $N(v) \cup N(w) \cup N(x)-\{u\}$ such that they are of the same color pairwisly. Now by the pigeon role principle $, v, w$ or $x$ has in its neighbors, 4 of that 10 vertices. This means that $v, w$ or $x$ is a single vertex. Without loss of generality we suppose that $v$ is single. By eliminating the color $v$ and using Lemma 3.1, we can find a valid color for vertex $v$ such that $L_{u}=(2,1,1,1)$ and this case have been treated in case 3.1.2.

Case 3.1.4: $L_{u}=(4,1)$
Let $N(u)=\{v, w, x, y, z\}, c(v)=c(w)=$ $c(x)=c(y)=1, c(z)=2$ and $c(u)=9$ then $f(u)=\{3,4,5,6,7,8\}$. If we can't find a valid color from $f(u)$ for $u$, then we must have $C_{u}(1,3), C_{u}(1,4), C_{u}(1,5), C_{u}(1,6), C_{u}(1,7)$ and $C_{u}(1,8)$. If one of the vertices $v, w, x$ or $y$ is single, then by eliminating the color of this vertex and applying Lemma 3.1 for it we can find a valid color such that $L_{u}=(3,1,1)$ and this case was handled above. Now suppose that none of the vertices $v, w, x$ and $y$ are single. Since we have 6 critical cycles involving $u$ therefore there exist in $N(v) \cup N(w) \cup N(x) \cup N(y)-\{u\} 12$ vertices such that their colors are from the set $\{3,4,5,6,7,8\}$ and they are of the same color pairwisly. Whereas none of the vertices $\{v, w, x, y\}$ are single ,so the colors which assigned to the neighbors of one of them (consider $v$ ) are $\{3,4,5, \alpha, 9\}$. We have two cases for $\alpha$. Either $\alpha \in\{3,4,5\}$ or $\alpha=9$ (because $v$ isn't a single vertex). If $\alpha \in\{3,4,5\}$ (consider $\alpha=3$ ) then by using Lemma 3.4 for $v$, we have $c(v) \in\{2,6,7,8, \alpha\}$. We have two cases for $c(v)$. If $c(v) \in\{6,7,8, \alpha\}$ then $L_{u}=(3,1,1)$ and this case was treated. Now suppose $c(v)=2$ and 6,7 and 8 aren't valid colors for $v$. suppose that two vertices $s$ and $t$ are neighbors of $v$ such that $c(s)=c(t)=$ 3. By this assumption $\{6,7,8\} \subset S C N(t)$ and $\{6,7,8\} \subset S C N(s)$ (because we can't choose 6 and 7 and 8 for $v$ ). Now we recolor vertex $v$ with color 1 ,since we have $C_{u}(1,3)$,then in neighbors of vertex $s$ or $t$ (consider $t$ ) we have two vertices such that their colors are 1 (one of them is $v$ ). Since $3 \in f(t)$ and $3 \notin S C N(v)$ so we can recolor vertex $t$ with color 3 and this yields $L_{v}=(1,1,1,1,1)$,therefore we can recolor $v$ such that $L_{u}$ becomes $(3,1,1)$. If $\alpha=9$ then we can color vertex $v$ with color 6 and we have $L_{u}=(3,1,1)$ and this case was handled above.
Case 3.1.5: $L_{u}=(5)$
Let $N(u)=\{v, w, x, y, z\}$ and $c(v)=c(w)=$ $c(x)=c(y)=c(z)=1$ and $c(u)=$ 9 then $f(u)=\{2,3,4,5,6,7,8\}$. If we can't find a valid color from $f(u)$ for $u$, then we must have $C_{u}(1,2), C_{u}(1,3), C_{u}(1,4), C_{u}(1,5)$, $C_{u}(1,6), C_{u}(1,7)$ and $C_{u}(1,8)$. This means that we have at least 14 vertices in $N(N(u))-\{u\}$ such that their colors are from the set $\{2,3,4,5,6,7,8\}$ and they are of the same color pairwisly. We have two cases:

Case a: One of the neighbors of $u$ (consider $v$ ) contains 4 vertices from that 14 vertices as its neighbors.
In this case vertex $v$ is a single vertex and we can recolor it such that $L_{u}=(4,1)$.
Case b: None of the vertices in $N(u)$ contains 4 vertices from that 14 vertices as its
neighbors.
In this case there exist a vertex in $N(u)$ such that it contains 3 vertices from that 14 vertices as its neighbors. Suppose that this vertex is $v$. Without loss of generality we can assume $S C N(v)=\{2,3,4, \alpha\}$. If $\alpha \in\{2,3,4\}$ or $\alpha=9$ then we can recolor vertex $v$ such that $L_{u}=(4,1)$.
Case 3.1.6: $L_{u}=(2,2,1)$
Let $N(u)=\{v, w, x, y, z\}$ and $c(v)=c(w)=$ $1, c(x)=c(y)=2, c(z)=3$ and, then $f(u)=$ $\{4,5,6,7,8\}$. If there is no valid color in $f(u)$ for $u$ then we have some critical cycles. All critical cycles are of the two following types $C_{u}(1, \alpha), 4 \leq \alpha \leq 8$ or $C_{u}(2, \beta), 4 \leq \beta \leq 8$. We can consider two
in the similar way.
Case a: We have $C_{u}(1,4), C_{u}(1,5)$, $C_{u}(1,6)$ and $C_{u}(1,7)$
In this case the vertices $v$ and $w$ are single and we can find a valid color for vertex $v$ which is neither 2 nor 3 . After changing the color of $v$ with a new color we have $L_{u}=(2,1,1,1)$ and this case was handled above.
Case b: We have $C_{u}(1,4), C_{u}(1,5)$, $C_{u}(1,6)$ and $C_{u}(2,7), C_{u}(2,8)$.
By this assumptions we have $\operatorname{SCN}(v)=$ $\{4,5,6,9, \alpha\}$. We have some cases for $\alpha$. If $\alpha \notin\{4,5,6,9\}$, then the vertex $v$ is a single vertex and we can recolor it such that $L_{u}$ becomes $(2,1,1,1)$. If $\alpha=9$, we can recolor $v$ with color 7, then $L_{u}=$ $(2,1,1,1)$. Let $\alpha \in\{4,5,6\}$. Without loss of generality we can assume $\alpha=4$, then $L_{v}=(2,1,1,1)$, by using Lemma 3.4 for vertex $v$ we have $c(v) \in\{2,3,7,8, \alpha=4\}$. We have three possible cases such as $c(v) \in$ $\{7,8, \alpha=4\}$ or $c(v)=2$ or $c(v)=3$. If $c(v) \in\{7,8, \alpha=4\}$ then $L_{u}=(2,1,1,1)$ and this case was treated above. If $c(v)=2$ then $L_{u}(3,1,1)$ and this case was handled above.
Let $c(v)=3$ and 2,7 and 8 aren't valid colors for $v$. Suppose that two vertices $s$ and $t$ are neighbors of $v$ such that $c(s)=c(t)=$ 4. By this assumption $\{2,7,8\} \subset S C N(t)$ and $\{2,7,8\} \subset S C N(s)$ (because we can't choose 2 and 7 and 8 for $v$ ). Now we recolor vertex $v$ with color 1 , since we have $C_{u}(1,4)$, then in neighbors of vertex $s$ or $t$ (consider $t$ ) we have two vertices such that their colors are 1 (one of them is $v$ ). Since $3 \in f(t)$ and $3 \notin S C N(v)$ so we can recolor vertex $t$ with color 3 and this yields $L_{v}=(1,1,1,1,1)$ therefore we can recolor $v$ such that $L_{u}$ becomes $(2,1,1,1)$.

Case 3.1.7: $L_{u}=(3,2)$
Let $N(u)=\{v, w, x, y, z\}, c(v)=c(w)=1$ and $c(x)=c(y)=c(z)=2$, then $f(u)=$ $\{3,4,5,6,7,8\}$. If we can't find a valid color for $u$ from $f(u)$ then we have 6 critical cycles containing $u$. Each cycle needs two vertices from $N(N(u))$ -
$\{u\}$ of the same color. Therefore there exist at least 12 vertices in $N(N(u))-\{u\}$ such that their colors are from $\{3,4,5,6,7,8\}$ an they are of the same color pairwisly. We detail two possible cases:

Case a: There exist a vertex in $N(u)$ such that it contains 4 vertices from that 12 vertices as its neighbors.
This vertex is a single vertex and we can recolor it. If this vertex is $v$ or $w$ then $L_{u}$ becomes $(3,1,1)$ and if this vertex is $x$ or $y$ or $z$ then we have $L_{u}=(2,2,1)$.
Case b: It doesn't exist a vertex in $N(u)$ such that it contains 4 vertices from that 12 vertices as its neighbors.
In this case there is a vertex such that it contains 3 vertices from 12 vertices as its neighbors (consider $v$ or $x$ ). First suppose that this vertex is $v$. We can assume $S C N(v)=\{3,4,5,9, \alpha\}$ (other cases for $S C N(v)$ can be handle by similar way). If $\alpha \notin\{3,4,5,9\}$ then $v$ is single and we can recolor it such that $L_{u}=(3,1,1)$. If $\alpha \in\{3,4,5\}$ then $L_{v}=(2,1,1,1)$ and by applying Lemma 3.4 for $v$, to recolor it, we will have $L_{u}=(3,1,1)$ (if $c(v) \neq 2$ ) or $L_{u}=(4,1)$ (if $c(v)=2$ ). If $\alpha=9$ then we can assign color 6 to $v$ and $L_{u}=(3,1,1)$. Now suppose that $x$ contains 3 vertices from those 12 vertices as its neighbors. By this assumption, we have $\operatorname{SCN}(x)=$ $\{3,4,5,9, \alpha\}$, If $\alpha \notin\{3,4,5,9\}$ then $x$ is a single vertex an after recolor it, we will have $L_{u}=(2,2,1)$. If $\alpha \in\{3,4,5\}$ (let $\alpha=3$ ) then $L_{x}=(2,1,1,1)$ and by using Lemma 3.4 to recolor $x$, we have $c(x) \neq 1$ or $c(x)=1$. If $c(x) \neq 1$ then $L_{u}=(2,2,1)$. Now suppose that $c(x)=1$ and none of the colors in $f(x)$, isn't valid color for $x$. Let $s, t$ are two neighbors of $x$ such that their colors are $\alpha=3$ (consider $t$ ). Since $c(x)=1$ and we had $C_{u}(2,3)$, therefore $S C N(t)=\{2,6,7,8,1\}(6,7,8$ are in $f(x)$ but not valid, $c(x)=1$ ). In this case we recolor $x$ with color 2 and this action implies that $L_{t}=(2,1,1,1)$. Because $1 \in f(t)$ and $1 \notin S C N(x)$, we can assign color 1 to $t$ and obtain $L_{x}=(1,1,1,1,1)$. Finally we can recolor $x$ such that $L_{u}=(2,2,1)$.

## IV. Conclusion

In this paper, we have shown that any graph of maximum degree 5 can be acyclically colored with 8 color. As far as lower bounds are concerned. We know that $a\left(K_{6}\right)=6$ then for $F$ family of graphs with maximum degree 5 we have $a(F) \geq$ 6. Closing the gap between those two bounds is a challenging open problem. In particular, we strongly suspect that the upper bound of 8 is not tight.

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